

# Defences to Curb Online Password Guessing Attacks

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**Abstract:** In this digital world, where huge amount of information is available online, illegitimate access to sensitive information is on the increase. This information is accessed using online password guessing attacks like brute force and dictionary attacks. In this paper we depict the inadequacy of existing protocols and we propose the Password Guessing Resistant Protocol (PGRP) which can effectively prevent these attacks. The system is very stringent for attackers and at the same time is very user friendly for legitimate users. The system prevents cookie theft related issues as it uses IP addresses to track known and unknown machines. It also makes use of ATTs to conquer the guessing attacks

**Keywords:** Online password guessing attacks, brute force attacks, dictionary attacks, ATT, CAPTCHA

## I. INTRODUCTION

The number of online users in real world has increased. Protecting confidential information using password based authentication system has thus become the need of the hour. Password is used as a means for authentication because it is easy for legitimate user to remember.

But, online guessing attacks on password based systems like brute force attacks and dictionary attacks are increasing nowadays. If these password guessing attacks succeed, confidential information stored on password based authentication system become vulnerable to illegitimate access. [1][2]The two most common online password guessing attacks are Brute Force attacks and Dictionary attacks.

### A. Brute force attacks or Exhaustive key search

It is a trial and error method used to obtain information such as user password. In this attacks automated software is used to generate a large number of consecutive guesses using combination of various upper and lower case letters, special characters, numbers

### B. Dictionary Attacks

A Dictionary attack attempts to defeat an authentication mechanism by systematically entering each word in a dictionary of likely passwords.

Both these attacks succeed because most of the users use simple passwords that can be easily guessed. Thus increase in usability has negative impact on security. Thus it is important to prevent online password guessing attacks. In this paper we present a password guessing resistant protocol (PGRP) which provides a security- usability balance which means that, it provides security without hampering usability of the user. PGRP builds on two previous proposals Pinkas and Sander (often called PS) and Van Oorschot and Stubblebine (often called VS).

PGRP uses IP addresses to track known and unknown machines. In particular, to limit attackers in control of a large botnet, PGRP enforces ATTs after a few failed login attempts are made from unknown machines. On the other hand, PGRP allows a high number of failed attempts from known machines without answering any ATTs.

Also in case where legitimate user fails to provide correct password, the system provides new password after the legitimate user passes few ATT challenges, this new password is send to legitimate user. This paper thus describes PGRP based system that would prevent online password guessing attacks by using ATTs, tracing IP address and is also user friendly.

It also tries to provide new password safely to the legitimate user in case if the legitimate user fails to provide the right password. The system thus treats known and unknown machine differently thus increasing security as well as user friendliness at the same time which at present most password based authentication systems fail to provide

## II. EXISTING SYSTEM

There exist two login protocols that prevent online password guessing attacks These protocols can be explained as under:

### A. PINKAS and SANDER(PS) PROTOCOL

This protocol asks the users both legal as well as attackers to pass the ATT first and allows them to enter the username and password if the answer made is correct.[4] The improved version of PS stores browser cookies i.e. the information about the machine of users who had previously login successfully. The cookie is related to username of the last successful login attempts. Once the user requests the login server URL, user's browser sends the cookie back to server.

The protocol then requests the user to enter a username, password pair. If the pair is correct and a valid cookie is received from browser then user is granted access. If the pair is correct but no valid cookie is received then an ATT challenge must be answered before account access is granted.

1) *DISADVANTAGES of PS PROTOCOL:*

- valid users must also pass an ATT challenge for every login attempt thus affecting user convenience.
- Performance issues exist as the login server has to generate an ATT challenge for every login attempt.

*B. VAN OORSCHOT STUBBLEBINE (VS) PROTOCOL*

VS proposed modification to previous protocol [4][3]. It tracks failed login per username to impose ATT challenges. The protocol maintains a threshold for failed login attempts. Thus the ATT challenge would be given for incorrect pair depending upon this threshold value.

In addition to it, if credentials entered by the user is correct but the cookie is not valid then the user is asked whether the machine in use is trustworthy and if user uses it regularly. The cookie is stored in user's machine only if user responds yes to the question. This approach helps to reduce password guessing due to cookie theft

1) *DISADVANTAGES of VS PROTOCOL:*

- The legal user always face an ATT challenge once the threshold is exceeded. This feature enables attackers to cause a DOS attack by initiating failed login attempts greater than threshold for each targeted username, forcing ATT challenges for subsequent login attempts
- Does not restrict the number of failed login attempts for attackers.

### III. PROPOSED SYSTEM

The proposed system named PGRP is built upon the existing PS and VS protocols.

#### A. WORKING

PGRP maintains a white list (W) to distinguish between known and unknown machine. White list maintains a list of pair of source IP address and user name. Known machines are the ones for which a successful login attempt was initiated using a valid username password pair from source IP address.

The rest are treated as unknown machines. PGRP maintains a list called FU which records the number of failed login attempts per username. PGRP maintains a list

called FP which records the number of failed login attempts per source IP, username pair. Where source IP address is the valid IP address in the White list or a host with valid cookie and username is valid username attempted from source IP address. Each entry in white list, FU, FP is valid for time interval t1, t2, t3 respectively. This can be implemented using time stamp.

The system also maintains two variables max1, max2 which decides the maximum number of login attempts for known machine and unknown machine respectively. max1 is always greater than max2 because legitimate users must be given more number of trial attempts. If the number of failed login attempts exceeds the threshold values max1 and max2 for known and unknown machines respectively then an Automated Turing Test (ATT) is flashed.

Use of ATTs helps prevent most of online guessing password attacks since these tests are generated by the computer but cannot be solved by it. Mostly CAPTCHAs are used as ATTs. If the answer to the ATT is correct user is granted access.

Except for one case wherein the username is valid but no valid cookie is received or host does not belong to white list. In that case the user is denied access even if answer to ATT is correct.

This feature prevent attacker from initiating brute force attack by restricting the attacker from knowing valid usernames. But if answer to ATT is incorrect then access is denied.

The proposed system distinguishes among the known and unknown machine and thus allows more privileges (More no. of trial attempts) to the known machine. It's a legitimate user friendly system. Figure.1 shows diagrammatic representation of the proposed PGRP system.

The general idea behind PGRP is that except for the following two cases, all remote hosts must correctly answer an ATT challenge prior to being informed whether access is granted or the login attempt is unsuccessful:

- 1) when the number of failed login attempts for a given username is very small; and
- 2) when the remote host has successfully logged in using the same username in the past (however, such a host must pass an ATT challenge if it generates more failed login attempts than a pre-specified threshold). In contrast to previous protocols, PGRP uses either IP addresses, cookies, or both to identify machines from which users have been successfully authenticated.

The decision to require an ATT challenge upon receiving incorrect credentials is based on the received cookie (if any) and/or the remote host's IP address. In addition, if the number of failed login attempts for a specific username is below a threshold, the user is not required to answer an ATT challenge even if the login attempt is from a new machine for the first time. Also after passing an ATT challenge a new password is generated for the username and this new password is being forwarded to the user's mobile

**B. PROPOSED SYSTEM DIAGRAM**

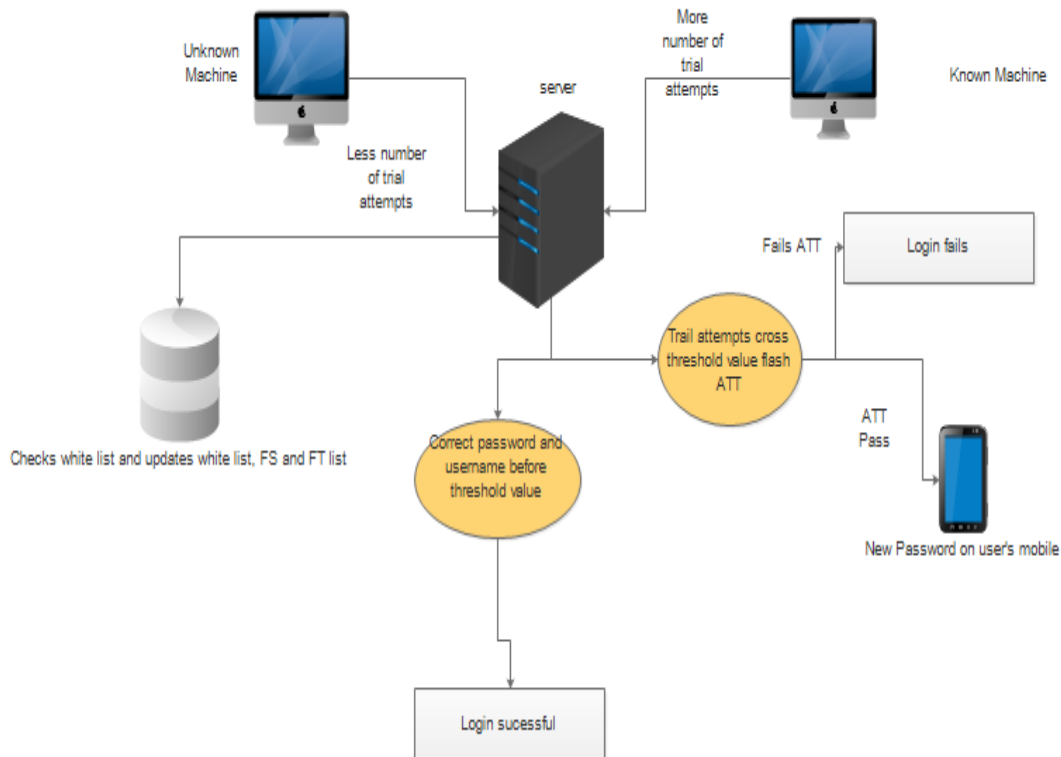


Fig. 1 Authentication system based on PGRP protocol

**IV. METHODOLOGY**

**Input:**

t1,t2,t3//time interval with default value t1=30 days,t2=1 day,t3=1 day  
w//whitelist which expires after time t1  
FU//number of failed login attempts per username  
FP//number of failed login attempts per username ,source IP pair  
max1//threshold value for known machine  
max2//threshold value for unknown machine.  
u//username  
p//password

ReadInput(u,p,cookie)

If LoginValid(u,p) then

    If(((valid(cookie,u,max1,true) or ((IP,u)belong to W))and(FP[IP,u]<max1))v(Fu[u]<max2)) then

        Fp[IP,u]=0

        Add IP to W

        GrantAccess(u, Cookie) and reset the password

    Else If ((Valid(cookie,u,max1,false) or ((IP,u) belongs to w)) and (FP[IP,u]<max1)) then

        FP[IP,u]=FP[IP,u]+1

        Alert("The username or password is incorrect")

    Else If(ATT()=Pass) then

        Alert(" username or password is incorrect")

    Else if(ValidUsername9u0and FU[u]<max2)) then

        FU=FU+1

        Alert("username or password is incorrect")

    Else

        Alert(" The answer to ATT is incorrect")

Each entry in W, FU, and FP has a “write-expiry” interval such that the entry is deleted when the given period of time (t1, t2, or t3) has lapsed since the last time the entry was inserted or modified. There are different ways to implement write-expiry intervals.

A simple approach is to store a timestamp of the insertion time with each entry such that the timestamp is updated whenever the entry is modified.

At anytime the entry is accessed, if the delta between the access time and the entry timestamp is greater than the data structure write-expiry interval (i.e., t1, t2, or t3), the entry

**D.ACTIVITY DIAGRAM**

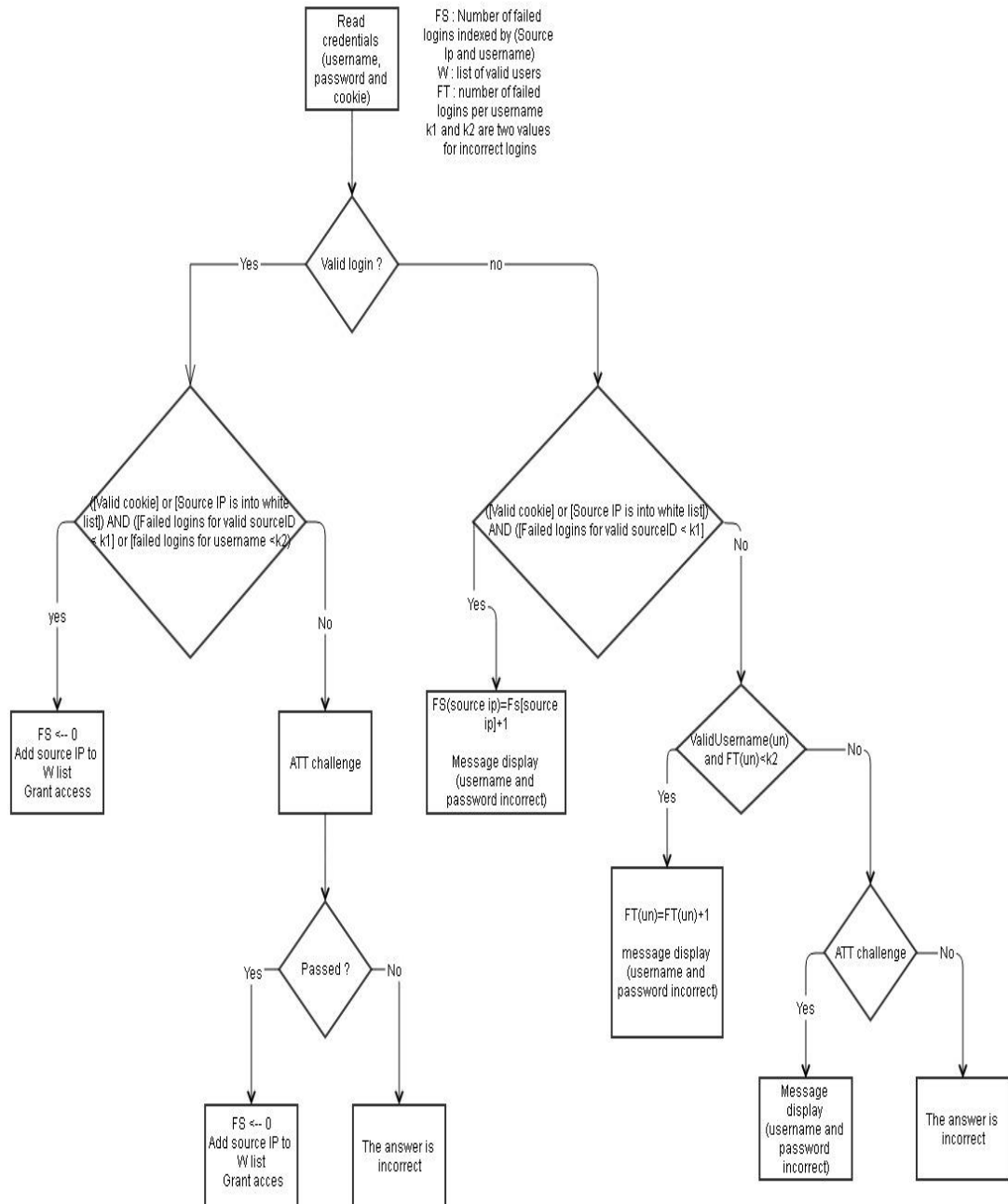


Fig. 2 Activity diagram of PGRP based system

**IV. COMPARISON of PGRP with EXISTING PROTOCOL**

TABLE 1

Sr no	Properties	PGRP	PS	VS
1	Limit the number of login attempts	Yes	No	No
2	Make brute force and dictionary attacks ineffective for large botnets .	yes	no	No
3	Impact on legitimate user convenience.	Less. As legitimate users are given more number of login attempts	Most. As legitimate users as well as attackers need to pass the ATT	More. As legitimate users undergo ATT only if threshold of failed login attempts is reached
4	Distinguish between known and unknown machine	yes	no	no
5	Graphical user interface supported	yes	yes	yes
6	Character user interface supported	yes	no	no
7	Issues related to cookie theft	No .As IP address is used	Yes. As cookies are used, it can be modified or deleted	Yes. As cookies are used it can be modified or deleted

The PGRP can be used in password based authentication system so as to prevent online password guessing attacks. [6]Modifications can be made to the proposed system to enhance its use in the future .In case of account locking the new password can be sent to legitimate users phone or email.

In addition an activity log may be displayed to every legitimate user on login. This feature would alert the legitimate user in case of attacks caused by attacker by logging in correctly from known machine .

Honeypots can be used capture information about the attacker. PGRP system provides stringent security against attackers but at the same time it is very user friendly for legitimate users.

It is better than the existing protocols as it provides good security without affecting user convenience .It can work on any operating system i.e. it is scalable .It is easy to implement i.e. it is deployable .

Online guessing attacks must be prevented because confidential information is always vulnerable to illegitimate access. PGRP is an excellent system that can prevent such attacks

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