



A Review to the Stability of Discrete Time State-Space Filters using Saturation Non-Linearity

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Abstract: The problem concerning the elimination of overflow oscillation in fixed-point state-space digital filter employing saturation arithmetic is considered by various researchers. In this paper a review is done to the finite procedure proposed by T. Oba [1] to test the stability of digital filters under saturation arithmetic.

Keywords: Digital filters, finite wordlength, asymptotic stability, nonlinear system.

I. INTRODUCTION

When a digital filter is implemented on a digital computer or on special-purpose digital hardware, the filter coefficients are stored in binary registers. These registers can accommodate only a finite number of bits and hence the filter coefficients have to be truncated or rounded-off in order to fit into these register. The finite-word length in recursive digital filter produces non-linearities, namely quantization and overflow. The presence of such non-linearities may result in the instability of the designed system. When dealing with the design and implementation of fixed-point state-space digital filters, it is, therefore, essential to know the conditions under which the filter will be globally asymptotically stable.

II. SYSTEM DESCRIPTION

The system under consideration is described by

$$x(r+1) = f(y(r)) = [f_1(y_1(r)) \quad f_2(y_2(r)) \quad \dots \quad f_n(y_n(r))]^T \quad (1a)$$

$$y(r) = [y_1(r) \quad y_2(r) \quad \dots \quad y_n(r)]^T = Ax(r) \quad (1b)$$

Where $x(r)$ is an n-vector space, $A = [a_{ij}]$ is the n x n coefficient matrix, and T denotes transpose. The saturation nonlinearity is given by

$$f_i(y_i(r)) = \begin{cases} 1 & y_i(r) > 1 \\ y_i(r) & |y_i(r)| \leq 1 \\ -1 & y_i(r) < -1 \end{cases} \quad (1c)$$

$i=1, 2, 3, \dots, n$ are under consideration.

Eq(1) is used to describe digital filters with symmetric saturation implemented with finite register length under zero external inputs.

III. THEOREM 1

The system described in (1) is asymptotically stable if there exists a positive definite matrix P satisfying

$$(P)_{i,i} - \sum_{j \neq i} (w_{|A|})_j (|P|)_{i,j} > 0 \quad \text{for all } i \in J_{|A|}^c \quad (2a)$$

such that $P - A^T P A$ is positive definite.

There are some prerequisite which are to be known before stating the algorithm to calculate $w_{|A|}$ for (2a), they are

- Stability test is to be done on matrix A, where $A \in R^{n \times n}$,
- The order of the matrix A is n.



- c) The matrix $B = [a_{ij}]$, $i, j = 1, 2, \dots, n$
- d) $J_0 = \phi$; J_k contains coordinate indices
- e) $n_0 = 0$; n_k contains the number of indices of J_k
- f) J_k^c contains the complement indices of J_k

g)

$$w_0 = 1 \in R^n ; \quad w_0 = \begin{bmatrix} 1 \\ 1 \\ 1 \\ \cdot \\ \cdot \\ \cdot \\ \cdot \\ n \end{bmatrix}$$

IV. ALGORITHM 1

The following procedure is proposed by ref [1] with $k=1$, to obtain J_B , n_B and $w_B = R^n$

- i) Let J_k denotes the list of coordinate indices i 's satisfying $(Bw_{k-1})_i < 1$, and let n_k denotes the number of the indices in J_k
- ii) If $J_k = J_{k-1}$, or if $n_k = n$, then define $J_B = J_k$, $n_B = n_k$, and

$$w_B = \begin{cases} w_{k-1} & \text{if } n_B < n \\ 0 & \text{if } n_B = n \end{cases} \quad (2b)$$

and then exit the loop.

- iii) Define $w_k \in R^n$ such that

$$\begin{cases} (w_k)_{J_k} = (I_{n_k} - B_{J_k, J_k})^{-1} \sum_{b=1}^{n-n_k} (B_{J_k, J_k^c} I_{n-n_k})_{a,b} & a = 1, 2, \dots, n_k \\ (w_k)_{J_k^c} = I_{n-n_k} \end{cases} \quad (2c)$$

and return to step (i) with $k=k+1$.

V. NUMERICAL EXAMPLE 1

To illustrate the algorithm for the stability test of fixed-point state-space digital filter with saturation arithmetic, a specific example of a third-order digital filter is considered with

$$A = \frac{1}{10} \begin{bmatrix} 1 & 5 & 4 \\ 10 & -2 & 5 \\ 0 & -3 & -1 \end{bmatrix}$$

According to the prerequisite of the algorithm Order of the matrix A is 3

$$B = \frac{1}{10} \begin{bmatrix} 1 & 5 & 4 \\ 10 & 2 & 5 \\ 0 & 3 & 1 \end{bmatrix}$$

$$J_0 = \phi$$

$$n_0 = 0, \text{ and}$$

$$w_0 = \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}$$



Iteration1

Step (i)

$$Bw_0 = \frac{1}{10} \begin{bmatrix} 1 & 5 & 4 \\ 10 & 2 & 5 \\ 0 & 3 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} = \begin{bmatrix} 1 \\ 1.7 \\ 0.4 \end{bmatrix}$$

$(Bw_0)_i < 1$, where 'i' is the indices values satisfying the given condition, in this step it is {3} Therefore $J_1 = \{3\}$ and $n_1=1$ (number of indexes in J_1)

Step (ii)

$$J_0 \neq J_1 \text{ and } n_1 \neq n$$

$$\text{Step(iii)} \quad \begin{cases} (w_1)_{J_1} = (I_{n_1} - B_{J_1, J_1})^{-1} \sum_{b=1}^{3-1} (B_{J_1, J_1^c} I_{n-n_1})_{a,b} \\ (w_k)_{J_1^c} = I_{n-n_1} \end{cases} \quad a = 1$$

Where J_1^c contains the complement indexes of J_1 i.e. $J_1^c = \{1,2\}$

Now

$$(w_1)_{J_1} = \left[1 - \frac{1}{10} \right]^{-1} \sum_{b=1}^2 \left(\begin{bmatrix} 0 & 3 \\ 10 & 1 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \right)_{1,b}$$

$$(w_1)_{J_1} = \left[\frac{9}{10} \right]^{-1} \sum_{b=1}^2 ([0 \quad 0.3])_{1,b}$$

$$(w_1)_{J_1} = \left[\frac{10}{9} \right] [0.3] = 0.3333$$

$$(w_1)_{J_1^c} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

Now return to step (i) of the algorithm, with $k=k+1$ i.e. $k=2$

Iteration 2

Step (i)

$$w_1 = \begin{bmatrix} 1 \\ 1 \\ 0.3333 \end{bmatrix}; \text{ since } J_1 = \{3\} \text{ therefore}$$

$$w_1(3) = (w_1)_{J_1}$$

$$Bw_1 = \frac{1}{10} \begin{bmatrix} 1 & 5 & 4 \\ 10 & 2 & 5 \\ 0 & 3 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 1 \\ 0.3333 \end{bmatrix} = \begin{bmatrix} 0.7333 \\ 1.3666 \\ 0.3333 \end{bmatrix}$$

$(Bw_1)_i < 1$, where 'i' is the indices values satisfying the given condition, in this step it is {1,3}

Therefore $J_2 = \{1,3\}$ and $n_2 = 2$ (the number of indices in J_2)

Step (ii)

$$J_1 \neq J_2 \text{ and } n_2 \neq n$$

Step (iii)

$$\begin{cases} (w_2)_{J_2} = (I_{n_2} - B_{J_2, J_2})^{-1} \sum_{b=1}^{3-2} (B_{J_2, J_2^c} I_{n-n_2})_{a,b} \\ (w_2)_{J_2^c} = I_{n-n_2} \end{cases} \quad a = 1,2$$



Where J_2^c contains the complement indexes of J_2 i.e. $J_2^c = \{2\}$

Now

$$(w_2)_{J_2} = \left[\begin{matrix} \left[\begin{matrix} 1 & 0 \\ 0 & 1 \end{matrix} \right] - \frac{1}{10} \left[\begin{matrix} 1 & 4 \\ 0 & 1 \end{matrix} \right] \end{matrix} \right]^{-1} \sum_{b=1}^1 \left(\frac{1}{10} \left[\begin{matrix} 5 \\ 3 \end{matrix} \right] [1] \right)_{a,b} \quad a = 1,2 \quad (w_2)_{J_2^c} = [1]$$

$$(w_2)_{J_2} = \left[\begin{matrix} 0.9 & -0.4 \\ 0 & 0.9 \end{matrix} \right]^{-1} \begin{pmatrix} 0.5 \\ 0.3 \end{pmatrix}_{a,1} \quad a = 1,2$$

$$(w_2)_{J_2} = \begin{bmatrix} 0.7037 \\ 0.0333 \end{bmatrix}$$

Return to step (i) of the algorithm with $k=k+1$, i. e $k=3$

Iteration 3

Step (i)

$$w_2 = \begin{bmatrix} 0.7037 \\ 1 \\ 0.0333 \end{bmatrix}; \text{ since } J_2 = \{1,3\} \text{ therefore}$$

$$Bw_2 = \frac{1}{10} \begin{bmatrix} 1 & 5 & 4 \\ 10 & 2 & 5 \\ 0 & 3 & 1 \end{bmatrix} \begin{bmatrix} 0.7037 \\ 1 \\ 0.0333 \end{bmatrix} = \begin{bmatrix} 0.7037 \\ 1.0703 \\ 0.3333 \end{bmatrix}$$

$$w_2 \begin{pmatrix} 1 \\ 3 \end{pmatrix} = (w_2)_{J_2}$$

$(Bw_2)_i < 1$, where 'i' is the indices values satisfying the given condition, in this step, it is again {1, 3}.Therefore

$J_3 = \{1,3\}$ and $n_3 = 2$ (the number of indices in J_3)

Step (ii)

$J_2 = J_3$, and $n_3 \neq n$

In step (ii) of iteration 3, one of the conditions stated in step (ii) of the algorithm is satisfied. Therefore we will define

$$J_B = J_3 = \{1,3\}$$

$n_B = n_3 = 2$ and

$w_B = w_2$, since $n_B < n$

$$(w_B) = (w_2) = \begin{bmatrix} 0.7037 \\ 1 \\ 0.3333 \end{bmatrix}$$

Exit the loop.

To calculate the value of P for the given A in Numerical Example 1, we will use MATLAB LMI tool box. The matrix P for given A in Numerical Example 1 comes out to be

$$P = \begin{bmatrix} 0.8951 & -0.0789 & 0.1553 \\ -0.0789 & 0.5423 & -0.0186 \\ 0.1553 & -0.0186 & 1.1414 \end{bmatrix}$$

Following the algorithm stated in IV, for the A given in numerical example 1 we have $J_{|A|} = \{1,3\}$ and $J_{|A|}^c = \{2\}$.

Considering $J_{|A|}$, $J_{|A|}^c$ and P, for the given A in Numerical example 1, we will check whether Theorem 1 is satisfied, i.e.

$$(P)_{i,i} - \sum_{j \neq i} (w_{|A|})_j (|P|)_{i,j} > 0 \quad \text{for all } i \in J_{|A|}^c \quad (3a)$$

In our case

$$(P)_{i,i} - \sum_{j \neq i} (w_{|A|})_j (|P|)_{i,j} > 0 \quad \text{for all } i \in J_{|A|}^c = \{2\} \quad (3b)$$



$$(P)_{2,2} - (w_{|A|})_1 (|P|)_{2,1} - (w_{|A|})_3 (|P|)_{2,3} \quad (3c)$$

$$(0.5423) - (0.7037) * (|-0.0789|) - (0.3333) * (|-0.0186|) = 0.4806$$

Thus the value of (3c) comes out to be greater than zero. Hence the system considered in the numerical example 1 is judged to be asymptotically stable according to Theorem 1. The same can also be verified by plotting the state trajectories of the numerical example 1. The figure 1 shows that the system under consideration is stable, as the next state of the system reaches zero with increasing iterations i.e. the output reaches zero with zero input

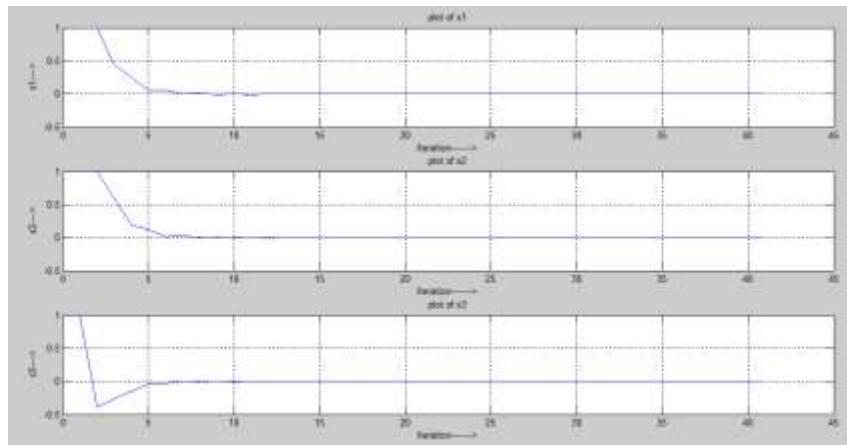


Figure 1 Dynamical behavior of the system considered in numerical example

VI. CONCLUSION

The criteria for the global asymptotic stability of fixed-point state-space digital filters with saturation nonlinearity have been given by several researchers. A finite procedure proposed by Ooba.T [1] ascertains the global asymptotic stability of the system considered in the numerical example. Modification is done to the algorithm proposed by [1], which is reasonably required and it broadens the scope of stability test from those of earlier results.

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