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International Journal of Advanced Research in Computer and Communication Engineering ISO 3297:2007 Certified Vol. 5, Issue 9, September 2016

File Replication for Achieving Querying Delay in Mobile Ad Hoc Networks Though Replication for Efficient File Sharing

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Abstract: Mobile computing is becoming more and more popular. The proficiency of file querying suffers from the properties of networks which include node mobility and limited communication range and resource. File sharing is one of the aspects which include peer to peer file sharing over MANET. Main advantages of P2P file sharing are files can be shared without base stations, overload on server can be avoided and it can exploit the otherwise. Wasted peer communication opportunities among mobile nodes. File replication which plays important role in enhancing file availability and reduce file querying delay. By creating replicas the probability of encountered requests can be improved. Random Way Point used for the normal MANET and Community-Based Mobility Model used for Disconnected MANETs. In RWP, nodes are moving with random speed to the randomly selected points, so the probability of meeting each node is similar for all the nodes Community-based mobility model used in some content dissemination or routing algorithms for disconnected MANETs. So both models contain idea of resource for file replication, which considers both node storage and meeting frequency.

Keywords: MANET, peer to peer network, File sharing, cooperative cache.

I. INTRODUCTION

In a mobile ad hoc network (MANET), mobile hosts can mobility and limited communication range and resource. communicate directly with one another using direct pair Existing ad hoc network mainly focused on global optimal wireless links. Because it requires no fixed infrastructure replica creation with minimum average querying delay. and most of the time no explicit administration a MANET New concept of resource for which considers both node can extremely useful to support communication in storage and meeting frequency. File replication is an challenging situations, such as in rural, remote, or disaster-struck areas. P2P computing refers to technology that enables two or more peers to collaborate spontaneously in a network of equals (peers) by using individual node replicates files it frequently queries or a appropriate information and communication systems without the necessity for central coordination. P2P networks are overlay networks typically operated on infrastructure (wired) networks, such as the Internet. However, the P2P overlay network is dynamic, where peers come and go (i.e., leave and join the group) for sharing files and data through direct exchange. Such peerto-peer communication paradigm will be very important in wireless multi-hop networks as centralized servers might not be available or located in the Internet. Therefore, P2P will be an interesting alternative for decentralizing services or making its own local resources available in the multi-hop network to serve local user communities. P2P overlay networks in the Internet and mobile ad-hoc networks share many key characteristics such as selforganization and decentralization due to the common nature of their distributed components.

II .EXISTING SYSTEM

The efficiency of file querying suffers from the distinctive • Cache miss properties of mobile ad hoc networks including node

effective way to enhance file availability and reduce file querying delay. It creates replicas for a file to improve its probability of being encountered by request. Each group of nodes create one replica for each file they frequently query. Unfortunately, it is impractical and inefficient to enable every node to hold the replicas of all files in the system considering limited node resource. Though redundant replicas are reduced by group cooperation, neighboring nodes may separate from each other due to node mobility, leading to large query delay.

Limitations on Existing System

- □ Leads to large query delay due to node mobility.
- Each individual node creates one replica it frequently queries, which cause waste of resources.
- Mobility of node affects the availability of files or message.

III.PROPOSED SYSTEM

File sharing in peer to peer mobile ad hoc network is made efficient and reduced delay with help increased cache size. In file sharing, it looks for two operations,

- Cache penalty

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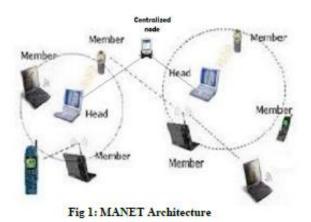
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piece of data in the cache, which results main memory matrix. access takes longer latency. Cache misses can be of three types: data read miss, instruction read miss and data write Node Insertion miss. Design and implementation of secured cooperative Each cell can have multiple numbers of nodes. Node count cache in wireless P2P networks are presented. Through for each cell has to be specified initially. Nodes are real implementations, important design issues are specified in a position with x, y points. Nodes in a network identified and proposed an asymmetric approach in order are movable since it is a mobile ad hoc network. to reduce the overhead of copying data between the user space and the kernel space, and also to reduce the data Gateway selection processing delay. The proposed algorithm well considers Each cell have multiple number of nodes within that the caching overhead and adapts the cache node selection gateway node is selected. Gateway is a router or a proxy strategy to maximize the caching benefit on different server that routes between networks. In this, gateway node MAC layers. Results show that the asymmetric approach acts as router between the cells. Always node which is outperforms the symmetric approach in traditional 802.11based ad hoc networks due to removal of most of the node in order to reach destination with reduced time and processing overhead.

IV. IMPLEMENTATION

A. System Design

Architectural diagram deals with cluster of mobile nodes. In this architecture each cluster contains cluster head. Each cluster head of each cluster is communicated with centralized node. Cluster head is also elected by centralized node. Each member in the cluster is identical in some features.



Each member in the cluster requests the cluster head for files and data. Cluster head includes all information about cluster members. Cluster head one group must communicate to cluster head of another cluster in order to get information of cluster member of that particular cluster.

B. Modules Description Cache and routing module

Cell Creation:

Cells are created with default X and Y axis. Each cell has Here authenticity is verified for checking whether a unique Id and position of each cell is displayed with help respective gateway node or any other node is active. of x,y points. All cells have a constant size which can be Authenticity also verifies for correct routing path, if the declared previously. Each cell have adjacent cells and path does not reach the destination reports a unsuccessful particularly with maximum of six adjacent cells. result.

A cache miss refers to a failed attempt to read or write a Adjacency of each cell is always specified in adjacency

near to the required adjacency cell is selected as a gateway delay.

Cache routing simulation module: Path finding:

Before finding the path between nodes. Source node and destination node are selected. With help of dynamic routing protocol DSR or AODV, efficient path can be established through gateway of each cell. Gateway could find exact direction of destination node. Use of dynamic routing protocol, whether DSR or AODV is previous selected. With that, distance, hope count, time duration can be found.

C. Caching and Routing module Authentication:

Diffie Hellman Digital Signature scheme is used as authentication technique in peer to peer mobile ad hoc network. Digital signature is mainly used for demonstrating the authenticity of the digital message or document. Digital signature is performed by a signing algorithm and it is verified by a verification algorithm

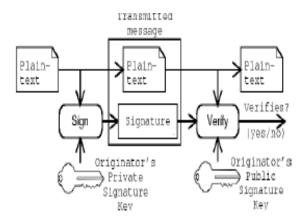


Fig 2: Digital Signature

ISSN (Online) 2278-1021 ISSN (Print) 2319 5940





International Journal of Advanced Research in Computer and Communication Engineering ISO 3297:2007 Certified

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Vol. 5, Issue 9, September 2016

V. IMPLEMANTATION

PATH		DSR		AODV		
DESTINATION NODE	HOP COUNT	DISTANCE	TIME	hop Count	DISTANCE	TIME
25	3	109	1.64	1	47	0.844
70	4	141	2.062	2	83	1.234
90	8	247	3.672	4	134	2.047
90	6	230	2.875	4	276	3.688
150	10	350	4.5	9	293	4.094
170	19	623	8.141	11	414	4.891
	DESTINATION NODE 23 70 90 90 150 170	DESTINATION NODE HOP COUNT 25 3 70 4 90 8 90 6 150 10 170 19	DESTINATION NODE HOP COUNT DISTANCE 25 3 109 70 4 141 90 8 247 90 6 230 150 10 350 170 19 623	DESTINATION NODE HOP COUNT DISTANCE TIME 25 3 109 1.64 70 4 141 2.062 90 8 247 3.672 90 6 230 2.875 150 10 350 4.5 170 19 623 8.141	DESTINATION NODE HOP COUNT DISTANCE TIME HOP COUNT 25 3 109 1.64 1 70 4 141 2.062 2 90 8 247 3.672 4 90 6 230 2.875 4 150 10 350 4.5 9 170 19 623 8.141 11	DESTINATION NODEHOP COUNTDISTANCE DISTANCETIME COUNTHOP COUNTDISTANCE2531091.641477041412.0622839082473.67241349062302.8754276150103504.5923

Fig 3: Analysis of DSR and AODV

Analysis Report of distributed routing protocol DSR and AODV represented in a tabular format. Report deals with different ranges of source and destination node values. Ranges will be of (1-50),(1-100),(1-150),(1-200). For each certain sample values are taken. Example for 1-50 range ^[7] source node value is 10 and destination node value is 25 which are between range1-50. Likewise same procedure repeated for other ranges and analysis of time, distance, hop count is made. Variation in values corresponding to Time taken, Distance, and hop count is noted. According to that simple graphical representation is made with respect to distance and time

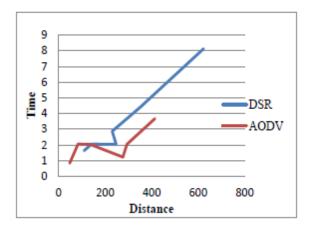


Fig 4: Distance vs Time

Graph discuss with analysis of distributed routing protocol DSR and AODV. The graph deals with the comparison of routing protocol for same source and destination values. Time taken, distance, hop count values are gradually high for DSR comparing to AODV routing protocol. Analysis Result is shown in the Graph. Analysis of distributed routing protocol DSR and AODV in p2p mobile ad hoc network is analyzed in implementation test bed Net Beans IDE 8.0.1. And final output representation is shown.

VI. CONCLUSION

Dynamic routing protocol DSR and AODV performance is enhanced in peer to peer mobile ad hoc network using gateway node and cooperative caching. Authenticated routing path is established from source to destination in peer to peer mobile ad hoc network with help of Diffie Hellman digital signature scheme.

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