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# Singly Linked List

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**Abstract**: This document gives formatting instructions for authors preparing papers for publication in the Proceedings of an International Journal. The authors must follow the instructions given in the document for the papers to be published. You can use this document as both an instruction set and as a template into which you can type your own text.

Keywords: nodes, address, data, link, head.

#### I. INTRODUCTION

Linked List can be defined as collection of objects called nodes that are randomly stored in the memory.

A node contains two fields i.e. data stored at that particular address and the pointer which contains the address of the next node in the memory. The last node of the list contains pointer to the null.

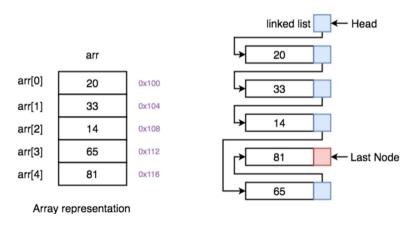


II. USES OF LINKED LIST

#### **Uses of Linked List**

- The list is not required to be contiguously present in the memory. The node can reside anywhere in the memory and linked together to make a list. This achieves optimized utilization of space.
- List size is limited to the memory size and doesn't need to be declared in advance.
- Empty node cannot be present in the linked list.
- We can store values of primitive types or objects in the singly linked list.

Below we have a pictorial representation showing how consecutive memory locations are allocated for array, while in case of linked list random memory locations are assigned to nodes, but each node is connected to its next node using pointer.





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#### Why use linked list over array?

Array contains following limitations:

- The size of array must be known in advance before using it in the program.
- Increasing size of the array is a time taking process. It is almost impossible to expand the size of the array at run time.
- All the elements in the array need to be contiguously stored in the memory. Inserting any element in the array needs shifting of all its predecessors.
- Linked list is the data structure which can overcome all the limitations of an array. Using linked list is useful because,

1. It allocates the memory dynamically. All the nodes of linked list are non-contiguously stored in the memory and linked together with the help of pointers.

2. Sizing is no longer a problem since we do not need to define its size at the time of declaration. List grows as per the program's demand and limited to the available memory space.

#### III. TYPES OF LINKED LIST

#### Types of Linked list are as follows-

1. Singly linked list

- 2. Doubly Linked List
- 3. Circular Linked List

#### Singly linked list

- Singly linked list can be defined as the collection of ordered set of elements. The number of elements may vary according to need of the program. A node in the singly linked list consist of two parts: data part and link part. Data part of the node stores actual information that is to be represented by the node while the link part of the node stores the address of its immediate successor.
- One way chain or singly linked list can be traversed only in one direction. In other words, we can say that each node contains only next pointer, therefore we cannot traverse the list in the reverse direction.
- Consider an example where the marks obtained by the student in three subjects are stored in a linked list as shown in the figure.



#### **Operations on Singly Linked List**

#### 1. Node Creation

```
struct node
{
    int data;
    struct node *next;
};
struct node *head, *ptr;
ptr = (struct node *)malloc(sizeof(struct node *));
```

#### 2. Insertion

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- A node can be added in three ways:
- Insertion at the beginning of the list
- Insertion at the end of the list
- Insertion in between the nodes

#### Algorithm to insert an element at the beginning of linked list:

1. Start

2. Create the node pointer \*temp

Struct node \* temp

- 3. Allocate address to temp using malloc
- temp = malloc(sizeof(struct node));

4. Check whether temp is null, if null then

Display "Overflow"

Else

```
temp-> info=data
```

temp-> next=start

5. Start=temp

6. stop

#### Algorithm to insert an element at the end of linked list:

```
1. Start
```

2. Create two node pointers \*temp, \*q

struct node \* temp, \*q;

3. q= start

4. Allocate address to temp using malloc

temp = malloc(sizeof(struct node));

5. Check whether temp is null, if null then

Display "Overflow"

else

temp-> info=data

temp-> next=null

6. While(q->next!=null)

q = q -> next

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7. q->next= temp

8. stop

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#### 3. Deletion

- i. Delete first node
- ii. Delete last node
- iii. Delete the in between node

<ul><li>(i) Deletion from the Beginning:</li><li>Suppose we have four nodes in the list and we want to delete</li></ul>	the first node.
start $\rightarrow$ 10 $\rightarrow$ 20 $\rightarrow$ 30 $\rightarrow$ 40	Niber (
Fig. 5.38	1910
• Now set, q = start	and the second second
start $\rightarrow$ 10 $\rightarrow$ 20 $\rightarrow$ 30 $\rightarrow$ 40 Fig. 5.39	
<ul> <li>Now make start to points to the node where starts ptr field i following statement.</li> </ul>	
start = start $\rightarrow$ ptr	ifisie
q → 10 • Start 20 • 30 • 40 P Sischwigbon add college Fig. 5.40	Now churge the ptr field is point
Now make q's ptr field to point to null and then free it statement. $q \rightarrow ptr = null$ free (q)	with the following
$q \rightarrow 10 \qquad $	40
Fig. 5.41 and the share of the Fig. 5.41 and the plant at a	Now change th
After deleting the node the list will look like.	painting.
$start$ 20 $\bullet$ 30 $\bullet$ 40 $\bullet$	tento → Now change th
Fig. 5.42	This will estable

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	Fig. 3.44
(ii)Dele	tion from the End:
• Trav	erse upto the second last node in the list set it as q.
	nge the ptr field of node q to the null.
Star Street	
	q temp
	start $\rightarrow$ 10 $\bullet$ 20 $\bullet$ 30 $\bullet$ 40 $\bullet$ 30 $\bullet$
	making the appropriate ling.
	There are three different places from 5.43
With	statement: Statement:
with	$q \rightarrow ptr = NULL$ are the tail of introduced to be tail of introduced to the second se
New	
	free the temp node free (temp).
After	deleting the last node the list will look like:
	start $\rightarrow$ 10 $\rightarrow$ 20 $\rightarrow$ 30 $\rightarrow$
	Now set, or a stand of the second contract war and the work on the
	and a provide the second
10-10-10-10-10-10-10-10-10-10-10-10-10-1	00 - 01 Fig. 5.44 00 - 01 - 01 - 031a
	start $\rightarrow$ 10 $\rightarrow$ 20 $\rightarrow$ 30 $\rightarrow$ 40 $\rightarrow$ 20 $\rightarrow$
	Fig. 5.45, at 1888
<ul> <li>Nov ptr</li> </ul>	v change the ptr field of the node one before q to point to the node where q's field is pointing.
• Afte	r that make q's ptr field to point to null with the following statements.
ginning	
node is	start > 10 = 20 = 30 = /> 40 =
	$r \rightarrow ptr = q \rightarrow ptr$
	$q \rightarrow ptr = null$
	Fig. 5.46
<ul> <li>Now</li> </ul>	free the node q free (q). After which the list will look like.
	start $\rightarrow$ 10 $\bullet$ 20 $\bullet$ 40 $\bullet$
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S. S. Alan	Fig. 5.47 pairs Jossie to neustramsigmi

#### **IV.CONCLUSION**

#### Advantages of Linked Lists

- They are a dynamic in nature which allocates the memory when required.
- Insertion and deletion operations can be easily implemented.
- Stacks and queues can be easily executed.
- Linked List reduces the access time.

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#### **Disadvantages of Linked Lists**

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- The memory is wasted as pointers require extra memory for storage.
- No element can be accessed randomly; it has to access each node sequentially.
- Reverse Traversing is difficult in linked list.

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